

Internet Engineering

241-461

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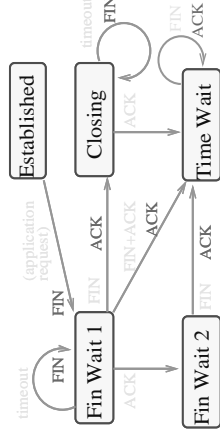
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TIME WAIT

◇ Why TIME WAIT ?



- Everything has happened?

- But that is at our system

 - What do we know about other system?

◇ Other system

- Received our FIN - ACKnowledged it

- Sent its own FIN - We ACKnowledged it

 - Did that ACK arrive?

 - How do we know?

TIME WAIT (2)

◇ Another issue

- Packets can be delayed

- Held in network - At a router usually

- Sometimes for a long time

 - Long by comparison with normal packets

- What happens if they arrive

 - After connection is CLOSED ?

◇ In CLOSED state we can open a new connection

- Same "connection" being used for new purpose

- Old data might be confused with new data ...

 - Sequence number & window help

 - But are not a guarantee

◇ TIME WAIT handles this

- Connection not CLOSED

 - So cannot be re-opened

- Wait in TIME WAIT long enough for old packets to die

Defining TCP

- ◇ Needs
 - Addressing
 - Src & Destination
 - Sequence Numbers
 - Acknowledge received data
 - Inform window size
 - Transmit Segments
 - Data accuracy protection
 - Magic Data
 - SYN & FIN
 - Segment maximum lifetime

Addressing

- ◇ Each Connection
 - My Network Address
 - My Port Number
 - Peer Network Address
 - Peer Port Number
- ◇ Network Addresses
 - Belong to Network Layer
 - Let Network layer handle those
- ◇ Port Numbers
 - Responsibility of Transport Protocol
 - TCP needs to handle them

Addressing (2)

- ◇ My vs Peer
 - One segment
 - Exists in Network
 - What is My ?
- ◇ Each segment is from one TCP to another
 - Source & Destination
- ◇ Each TCP segment needs
 - Source Port Identifier
 - Destination Port Identifier

Information Encoding

- ◇ Text vs Binary
 - Human vs Computer
- ◇ TCP uses BINARY
 - More efficient
 - Network bandwidth
 - Processing costs
- ◇ Binary encoding issues
 - Not all computers the same
 - Bigger problem in the past
 - Still exists today

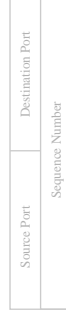
Port Identifiers

- ◇ Identifier just a number
 - An integer
 - 16 bits
 - convenient size
 - big enough
 - not too big
- ◇ TCP Header (so far)



Sequence Number

- ◇ 32 bit number
 - Needed in every segment



Acknowledgment

- ◇ 32 bit number
 - Sequence next expected
 - So must have same value range as SEQ

Source Port	Destination Port
Sequence Number	
Acknowledgment Number	
A	

- ◇ But not in every segment
 - Cannot send ACK before receiving SEQ
- ◇ So, just in most segments

Window Size

- ◇ 16 bit number
 - Return to size of window size later
- ◇ In every segment

Source Port	Destination Port
Sequence Number	
Acknowledgment Number	
A	Window Size

Data Reliability

- ◇ Checksum
 - Verifies data unmodified
 - Accidental modifications
 - ▷ not deliberate attacks

Source Port	Destination Port
Sequence Number	
Acknowledgment Number	
A	Window Size
Checksum	

Future Expansion

- ◇ Protocol might need revision
 - Or data not in every segment
- ◇ Add room for expansion
 - Or optional data

Source Port	Destination Port
Sequence Number	
Acknowledgment Number	
Hdr Len	A
	Window Size
Checksum	
Expansion Space	

- ◇ But how big is this?
 - Need size field

Magic Data

- ◇ Need to be able to send
 - SYN
 - FIN
- ◇ Those are data
 - but not normal data
- ◇ Only ever once in a segment (each)
 - Exist or do not exist

Source Port	Destination Port
Sequence Number	
Acknowledgment Number	
Hdr Len	A
	SYN
	Window Size
Checksum	

Complete TCP Header

Source Port	Destination Port						
Sequence Number							
Acknowledgment Number							
Hdr Len	Reserved	U	A	P	R	S	F
Checksum		Window Size		Urgent Pointer			

- ◇ Sent with every segment
 - Contains some data not examined yet

TCP Header

Source Port	Destination Port
Sequence Number	
Acknowledgment Number	
Header Reserved	U A P
Checksum	R S F
Window Size	
Urgent Pointer	

- ... Contains some data not examined yet

◊ Reset

- The TCP error message
- Sent whenever TCP error occurs
 - Segment in invalid state
 - NOT missing packet
 - NOT out of order packet
 - NOT duplicate packet

TCP Timeouts

◊ Recall ...

- retransmit
- after timeout

◊ How long?

- 1 second ?
- 2 seconds ?
- 3 seconds ?
- Long enough
 - no unnecessary retransmit
- Short enough
 - avoid delays

◊ Time to wait varies

- Different peers (near or distant)
- Different network state (idle or congested)

◊ Measure Round Trip Time

Next
step

Round Trip Time Estimation

◊ Useful to know when to retransmit

- if have not received ACK within the RTT
 - (plus a bit)
 - then assume packet lost

◊ But how to measure the RTT?

- Measure delay between packet and its ACK
 - easy

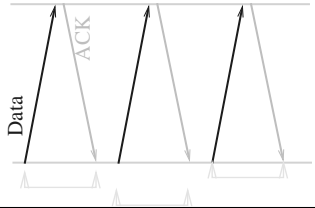
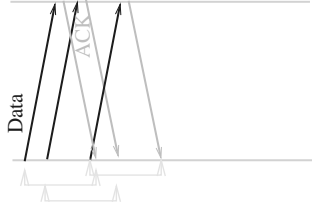
◦ But

- Send packet
 - wait ... wait ... wait (nothing)
- Retransmit packet
- ACK arrives

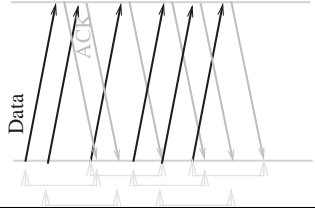
◦ Which packet was acknowledged?

- The initial packet
 - acknowledged slower than expected
- Or the retransmit?

TCP RTT Measurement



TCP RTT Measurement (2)



TCP RTT Measurement (3)

